

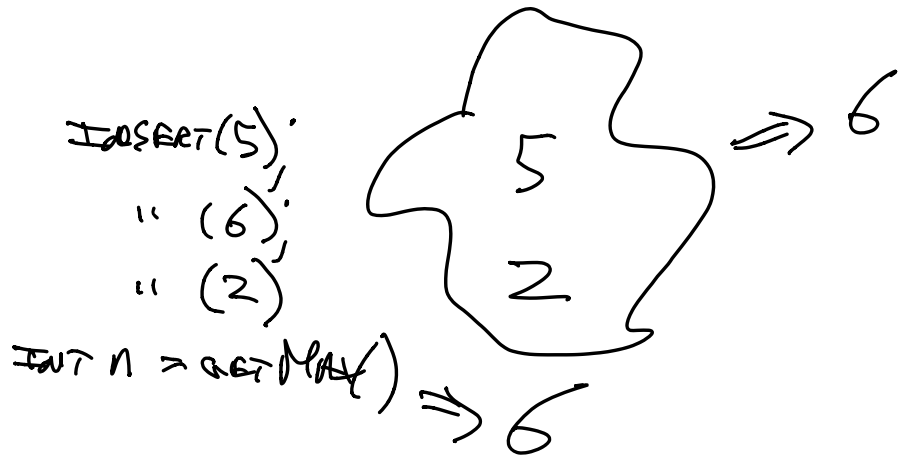
HEAP ("BINARY HEAP")

IS MOST IMPORTANT ~~ALGORITHM~~ ALGORITHM /
AS FAR PRIORITY QUEUES.

INTERFACE

```
PUBLIC INTERFACE INTEGER MAX QUEUE {  
    PUBLIC VOID INSERT(INT K);  
    PUBLIC PUBLIC INT GETMAX();  
    ... SIZE, ISEMPTY...  
}
```

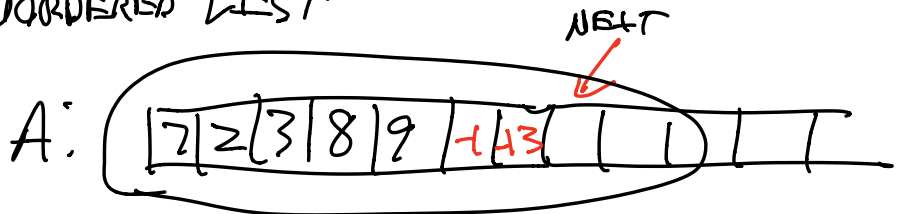
~



FIRST ATTEMPT

⇒ 12

UNORDERED LIST



$N = \# \text{ ELEMENTS}$
 NOT LENGTH OF ARRAY.

$O(1)$
 $O(N)$

INSERT: ADD TO END
 $A[\text{NEXT}++] = k$

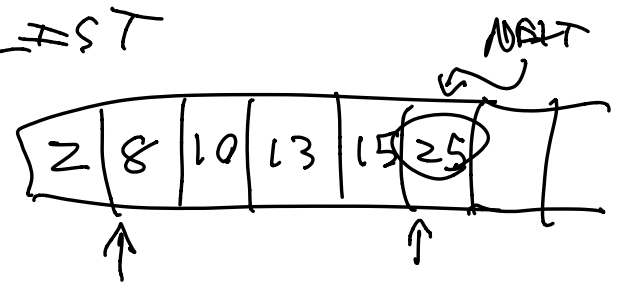
GET MAX: FIND MAX
 DELETE
 MAKE LATER IN
 QUER

$O(N)$

2ND ORDERED LIST

$O(N)$

INSERT

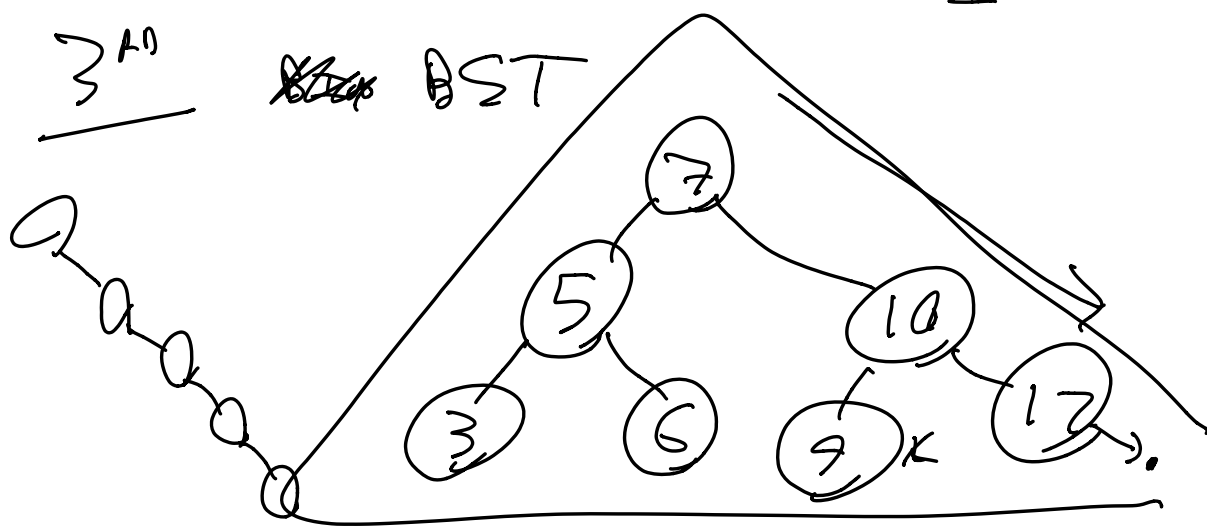


$O(1)$

GET MAX

RETURN $A[--\text{NEXT}]$

3RD ~~Step~~ BST



INSERT $O(N)$ } WORST
 GET MAX $O(N)$ } CASE

4TH

2-3 TREE

INSERT $O(\log(N))$

GET MAX $O(\log(N))$



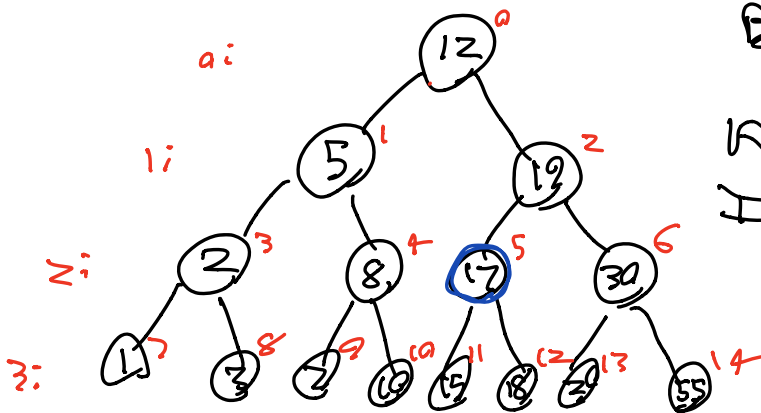
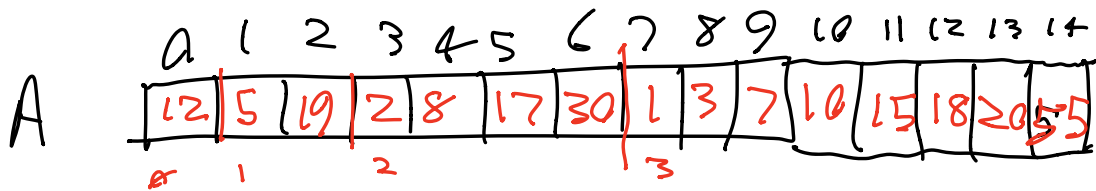
PUNCH LINE: WE WILL COMBINE
 BEST FEATURES OF
 2-4 ~~AND~~ OF THESE

HEAP DATA STRUCTURE

* ARRAY BASED (NO EXPLICIT
 PTRS)
 ORDERED
 TREE (BALANCED)

$O(\log(N))$ FOR ~~INSERT & GET MAX~~
 INSERT & GET MAX

STARTING A BINARY TREE IN AN ARRAY.



BASIC IDEA
 STORE KEYS IN TREE
 IN LEVEL ORDER

WE WILL REFER TO ITEMS BY INDEX
 $I = 5$ IN A

$$A[I] \Rightarrow 17$$

// RETURN INDEX OF PARENT OR -1 IF AT ROOT

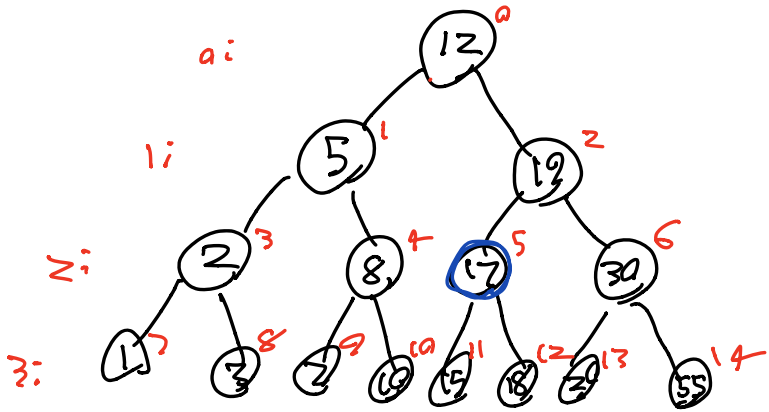
$$\text{PARENT}(5) \Rightarrow 2$$

INT PARENT(int I) {
 RETURN (I-1) // 2;
 }

| I | PARENT(I) | $I//2$ | $(I-1)//2$ |
|--------------|---------------|--------------|--------------|
| 0 | 1 | 0 | 0 |
| 1 | 0 ✓ | 0 | 0 |
| 2 | 0 | 1 | 0 |
| 3 | 1 | 1 | |
| 4 | 1 | 2 | |
| 5 | 2 | 2 | |
| 6 | 2 | 3 | |
| 7 | 3 | 3 | |
| 8 | 3 | 4 | |
| 9 | 4 | 4 | |
| 10 | 4 | | |
| 11 | 5 | | |
| 12 | 5 | | |

A

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |
|----|---|----|---|---|----|----|---|---|---|----|----|----|----|----|
| 12 | 5 | 19 | 2 | 8 | 17 | 30 | 1 | 3 | 7 | 10 | 15 | 18 | 20 | 55 |
| | 1 | | 2 | | | | 3 | | | | | | | |



```

INT LCHILD (INT I)
RETURN 2*I + 1

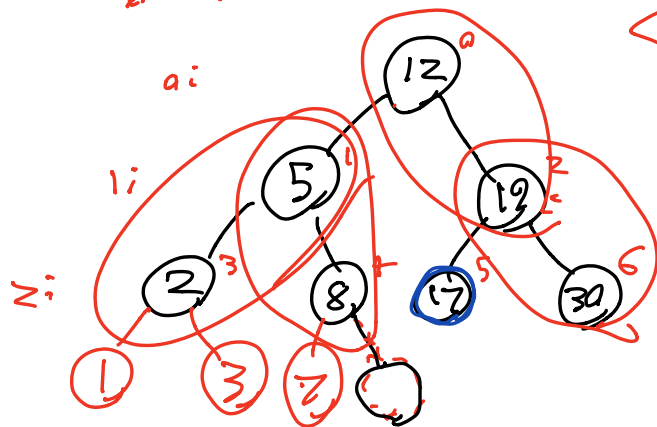
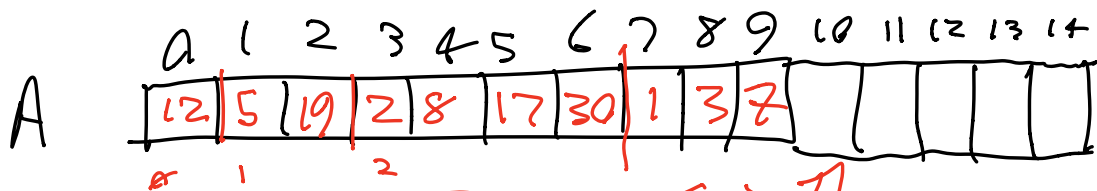
```

```

INT RCHILD (INT I)
RETURN 2*I + 2

```

| I | LC | RC | 2*I |
|---|----|----|-----|
| 0 | 1 | 2 | 0 |
| 1 | 3 | 4 | 2 |
| 2 | 5 | 6 | 4 |
| 3 | 7 | 8 | 6 |
| 4 | 9 | 10 | 8 |
| 5 | 11 | 12 | 10 |



$A[n_{left} + 1] = R$

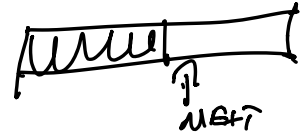
$\leftarrow n$

$\leftarrow n+1$

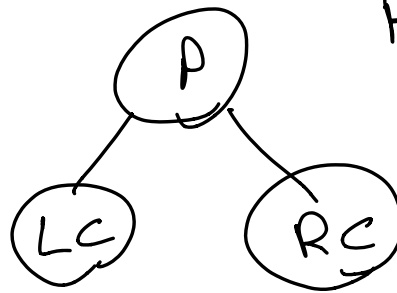
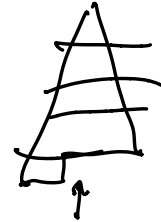
IF TREE IS "BALANCED FROM LEFT" IF BALANCED AND ALL LEAVES AT LEVEL n ARE TO RIGHT OF LEAVES AT LEVEL $n+1$.

FILL IN NODES IN LEVEL ORDER.

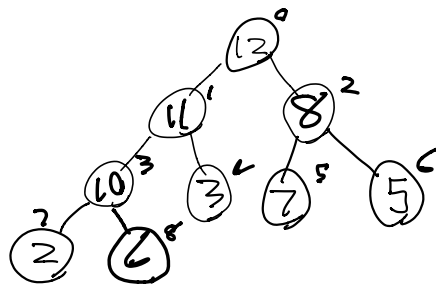
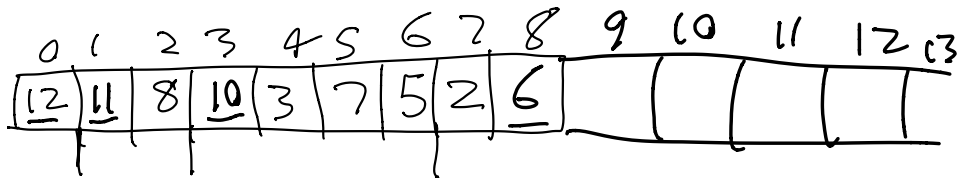
A HEAP IS AN ARRAY
 IN PREVIOUS FORM
 AND HAS "HEAP ORDERING
 PROPERTY" (ASSUMES MAX Q)



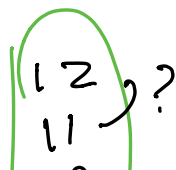
NAMEL: FOR ANY PARENT
 $P, P \geq LC$ & $P \geq RC$



$P \geq RC$



INSERT: PUT NEW KEY IN ^{MAX} LEVEL ORDER POSITION



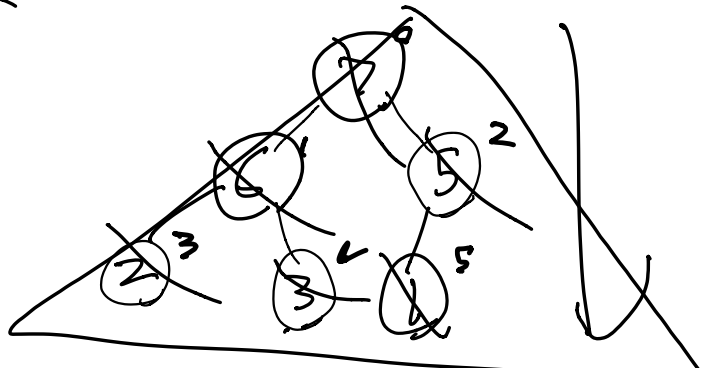
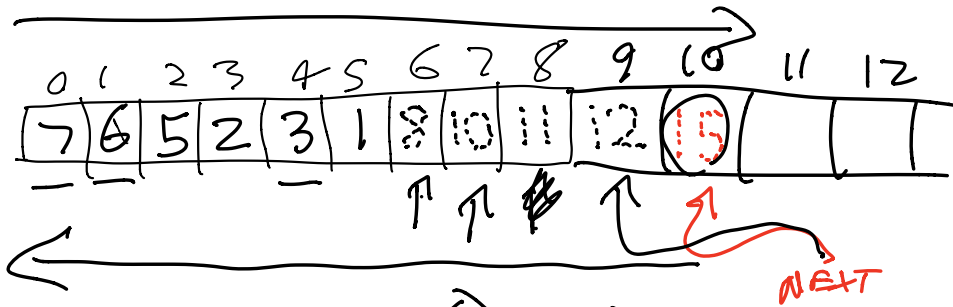
INSERT NEW KEY

10
6

IN ORDER IN
PATH TO ROOT
JUST AS IN
INSERTION SORT

COMPARE WITH
PARENT
SWAP IF
LARGER THAN
PARENT.

$N == \text{NEXT}$



-- NEXT;

BALANCED!
~~NEW~~ INSERT &
 GET MAX
 ARE $O(\log N)$
 USES ARRAY (CAN REVERSE)

DELETE MAX
 SWAP ROOT
 WITH LAST KEY
 SWAP NEW ROOT
 WITH MAXIMAL
 CHILD UNTIL
 IN HEAD ORDER.
 RETURN $A[\text{NEXT}]$

—

1